Non-uniform Hypercoherences

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Institut de Mathématiques de Luminy

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Denotational semantics (of λ -calculus, PCF, LL)

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Denotational semantics (of λ -calculus, PCF, LL) proofs, terms \rightarrow some math. structures: *agents* ٩ sequentiality & determinism. No concurrency ٩ **Internal interactivity** of calculus extensionality (PER), extensional collapse (quotient) ٩ uniformity: the stage where agents interact ٩ **Exponential** modalities of linear logic (! and ?)

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complexity issue (ELL, LLL), polarities, ...

 $P = \lambda \overline{b}. \quad \text{if } b \quad \text{then } \{ \text{ if } b \quad \text{then } t_1 \text{ else } t_2 \}$ else { if b then t_3 else t_4 }

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 $P: \mathsf{Bool} \to \mathsf{Bool} \cong !\mathsf{Bool} \multimap \mathsf{Bool}$ $t_i^{\bullet} \in \{t, f\}$

Coherence spaces and **hypercoherences** semantics (set-based) are *uniform*

$$P^{\bullet} = \{(\{\mathtt{t}\}, t_1^{\bullet}), (\{\mathtt{f}\}, t_4^{\bullet})\}$$

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 $extsf{P:Bool} o extsf{Bool} \cong extsf{Bool} o extsf{Bool}$ $t_i^{ullet} \in \{ extsf{t}, extsf{f}\}$

Relational semantics is *non-uniform*

 $P^{\bullet} = \{ ([\texttt{t},\texttt{t}], t_1^{\bullet}), ([\texttt{t},\texttt{f}], t_2^{\bullet}), ([\texttt{t},\texttt{f}], t_3^{\bullet}), ([\texttt{f},\texttt{f}], t_4^{\bullet}) \}$

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Multiset-based coherence spaces and hypercoherences semantics are also *uniform*

 $P^{\bullet} = \{ ([\mathtt{t}, \mathtt{t}], t_1^{\bullet}), ([\mathtt{f}, \mathtt{f}], t_4^{\bullet}) \}$

Static semantics

Dynamic semantics

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Static semantics: Agents are sets of points

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Static semantics: Agents are sets of points

Dynamic semantics: Agents are sets of pol. paths

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Static semantics: Agents are sets of **points** • additives \rightarrow disjoint unions, \emptyset

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- additives \rightarrow disjoint unions, \emptyset
- multiplicatives \rightarrow product of sets, $\{*\}$

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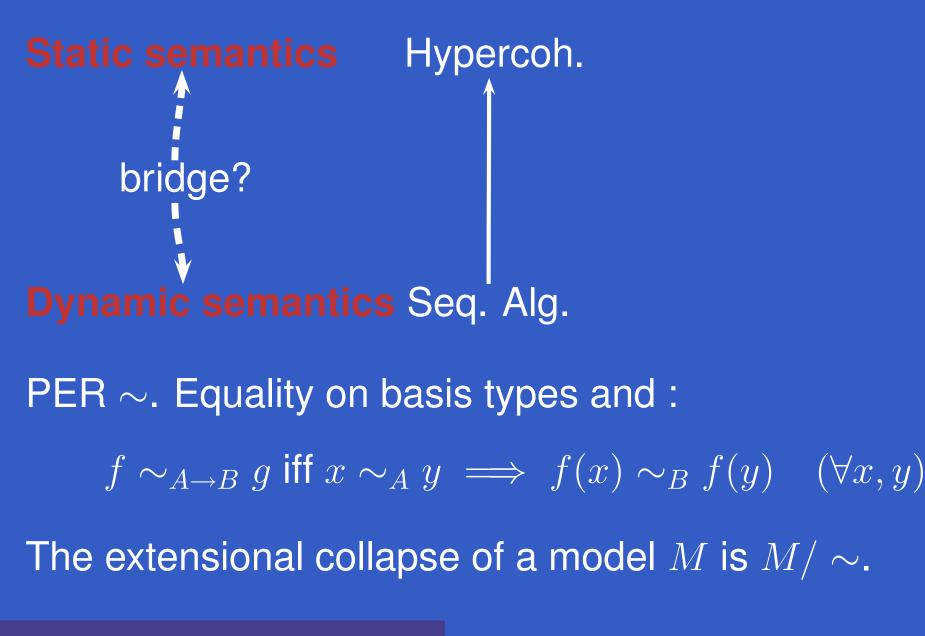
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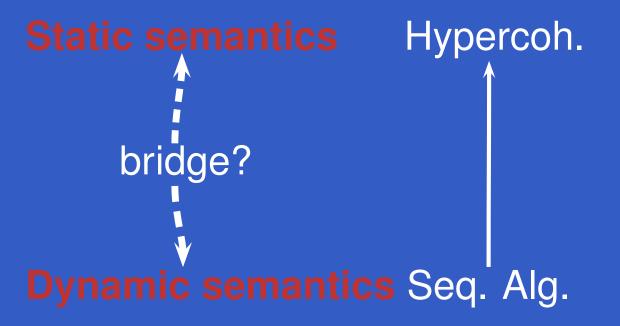
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- Dynamic semantics: Agents are sets of pol. paths
 - additives \rightarrow disjoint unions, \emptyset
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 - exponentials → interleaving of paths in agents uniformity

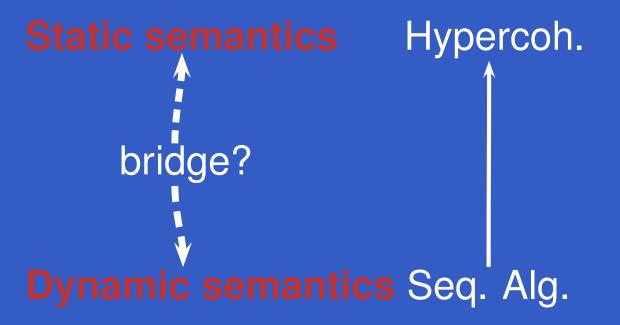
Static semantics bridge? Dynamic semantics

Static semantics Hypercoh. bridge? Dynamic semantics Seq. Alg.

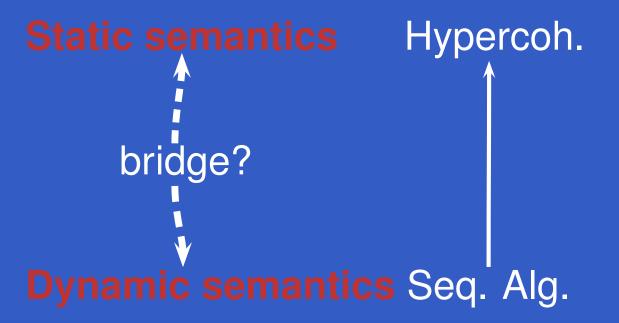




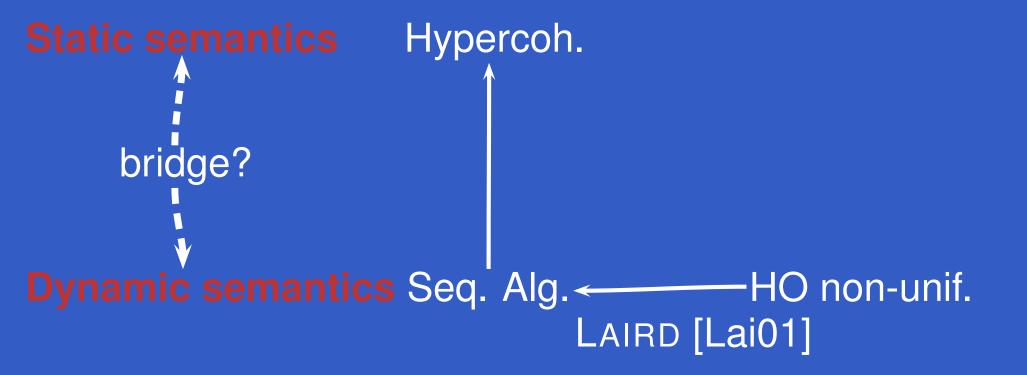
one-way bridge

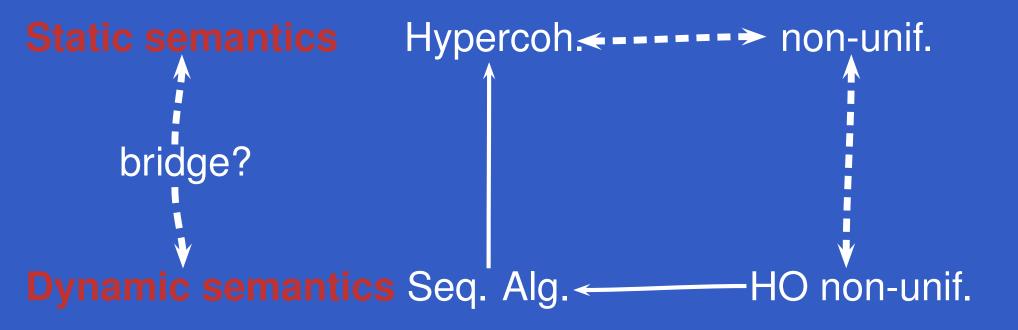


- one-way bridge
- abstract result



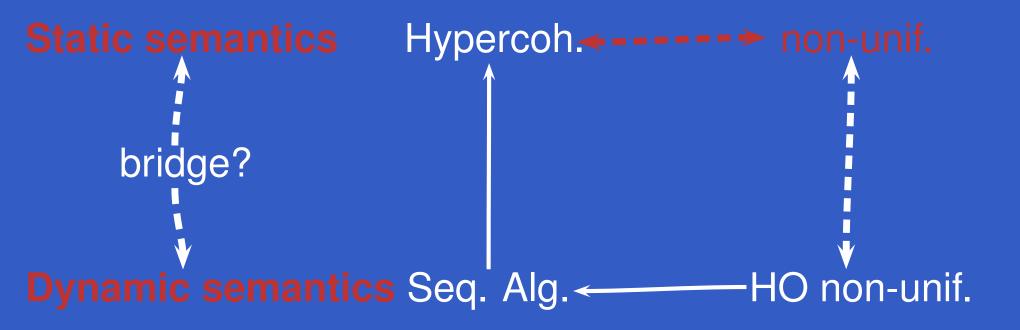
- one-way bridge
- abstract result
- Only simple types. Not the full power of linear logic.





Shifting to non-uniformity might improve the bridge

Denotational semantics



Shifting to non-uniformity might improve the bridge

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A **Power** *P* is a functor from the category of sets and inclusions to itself.

• Coherence space $X = (|X|, \bigcirc_X)$: graph $\{[a, a] \mid a \in |X|\} \subseteq \bigcirc_X \subseteq \mathcal{M}_{\{2\}}(|X|)$ • Hypercoherence $X = (|X|, \bigcirc_X)$: hypergraph $\{\{a\} \mid a \in |X|\} \subseteq \bigcirc_X \subseteq \mathcal{P}^*_{\text{fin}}(|X|)$ Reflexivity

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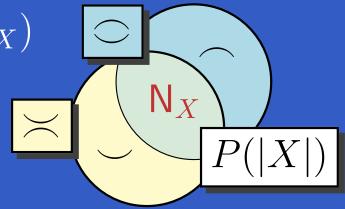
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P-coherence space $(|X|, \bigcirc_X, \asymp_X)$



P-coherence space $(|X|, \bigcirc_X, \asymp_X)$ No longer reflexive. Neutrality.

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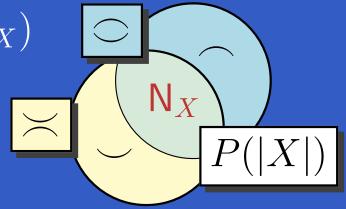
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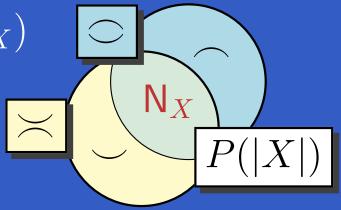
 \rightarrow a class of non-uniform coherence semantics which we describe now using the Power \mathcal{M}_K (multisets whose cardinalities are in K) where $K \subseteq \mathbb{N} \setminus \{0, 1\}$.

 \mathcal{M}_K -coherence space $(|X|, \bigcirc_X, \overleftarrow{\prec}_X)$



 \mathcal{M}_K -coherence space $(|X|, \bigcirc_X, \asymp_X)$

MALL model: standard pattern

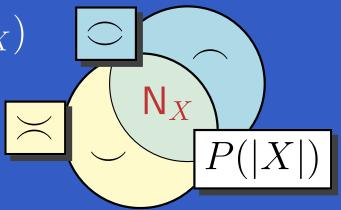


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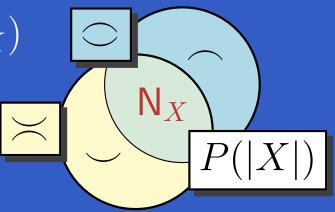


eg, $|X \multimap Y| = |X| \times |Y|$ and, for $s \in \mathcal{M}_K(|X \multimap Y|)$: $s \in \bigcirc_{X \multimap Y} \text{ iff } \begin{cases} \pi_1(s) \in \bigcirc_X \implies \pi_2(s) \in \bigcirc_Y \\ \pi_1(s) \in \asymp_X \Longleftarrow \pi_2(s) \in \asymp_Y \end{cases}$ $s \in N_{X \multimap Y}$ iff $\pi_1(s) \in N_X$ and $\pi_2(s) \in N_Y$

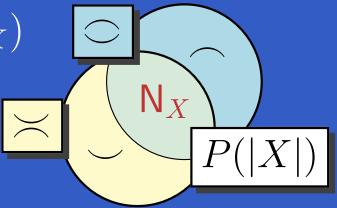
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- for MALL neutrality is reflexivity



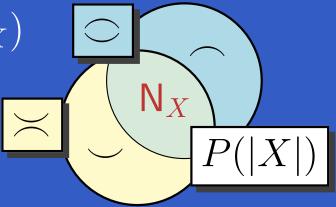
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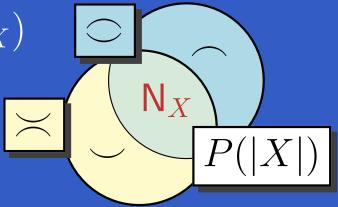
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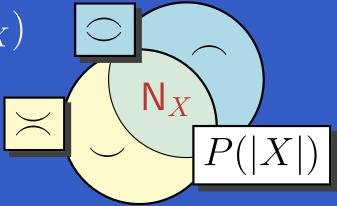
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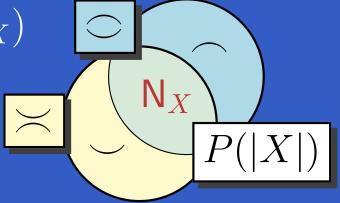
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- Hopefully, we have a better solution!



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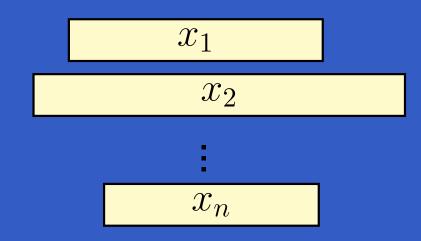
Web: $|X| = \mathcal{M}_{fin}(|X|)$

Web: $|!X| = \mathcal{M}_{fin}(|X|)$ This web is the same as for relational model, so the semantics is non-uniform. A uniform web would have been $\{x \in \mathcal{M}_{fin}(|X|) \mid \operatorname{supp}(x) \in \operatorname{agent}(X)\}$.

Web: $|!X| = \mathcal{M}_{fin}(|X|)$ Coherence. For each $[x_i \mid i \in I] \in \mathcal{M}_K(|!X|)$ we set:

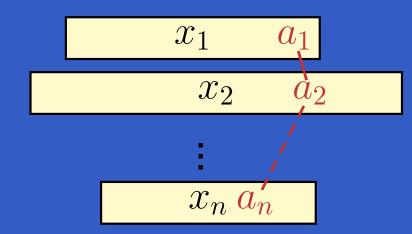
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• $[x_i \mid i \in I] \in \mathsf{N}_{!X}$ iff $[x_i \mid i \in I] \notin \smile_{!X}$ and $\exists (a_i^j)_{i \in I}^{j \in J}$,

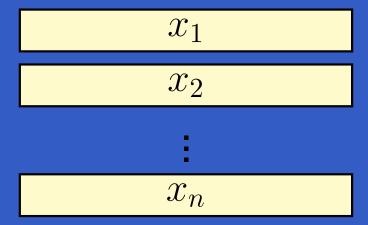
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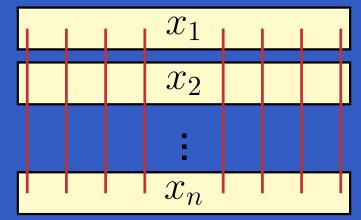
• $\forall i \in I, [a_i^j \mid j \in J] = x_i \text{ and }$



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- $[x_i \mid i \in I] \in \mathsf{N}_{!X}$ iff $[x_i \mid i \in I] \notin \smile_{!X}$ and $\exists (a_i^j)_{i \in I}^{j \in J}$,
 - $\forall i \in \overline{I, [a_i^j \mid j \in J]} = x_i$ and
 - $\forall j \in J, [a_i^j \mid i \in I] \in \mathsf{N}_X$





Non-uniform model of linear logic (logical forgetful functor).

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- $l_u = N_K!$. A new class of uniform models related to the non-uniform ones in a very comfortable way.

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• Forget multiplicities \rightarrow non-unif. hypercoherences ($!_{nuh} = S!$). Usual (multiset-based) uniform semantics. $!_{mh} = N !_{nuh}$.

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 Multicoherences ≠ hypercoherences → two (extensionally) different notions of higher order sequentiality. (Contrarily to what was expected, eg in [Lon02]).

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