

Projecting games on hypercoherences

polarized bordered games

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Games and hypercoherences

Denotational semantics:

if $\pi \rightarrow_* \pi'$ then $[\pi] = [\pi']$

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(two-players plays)

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results of interaction
(points/vertices)

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Comparison?

First answer: extensional collapse

Sequential algorithms

Hypercoherences



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Sequential algorithms

strategies
algorithm



extensional collapse
Ehrhard 1999



Hypercoherences

cliques
function

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strategies $\xrightarrow{\text{extensional collapse}}$ cliques
algorithm Ehrhard 1999 \longrightarrow *function*

Quotient by partial equivalence relations (Kreisel 50'):

$$x \sim_l x \quad \text{and} \quad f \sim_{\sigma \rightarrow \tau} g \text{ iff } x \sim_{\sigma} y \implies f(x) \sim_{\tau} g(y)$$

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- Better detailed by Melliès 2003.
- Limited to simple types ($\sigma := \iota \mid \sigma \rightarrow \sigma$).

Another bridge: direct projection



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Games strategies  time-forgetting projection **Hypercoherences** sets of points

in general, the image of a strategy is not a clique

Another bridge: direct projection

Games
strategies

time-forgetting
projection

Relational model
relations



relational model \approx hypercoherences without coherence

Another bridge: direct projection

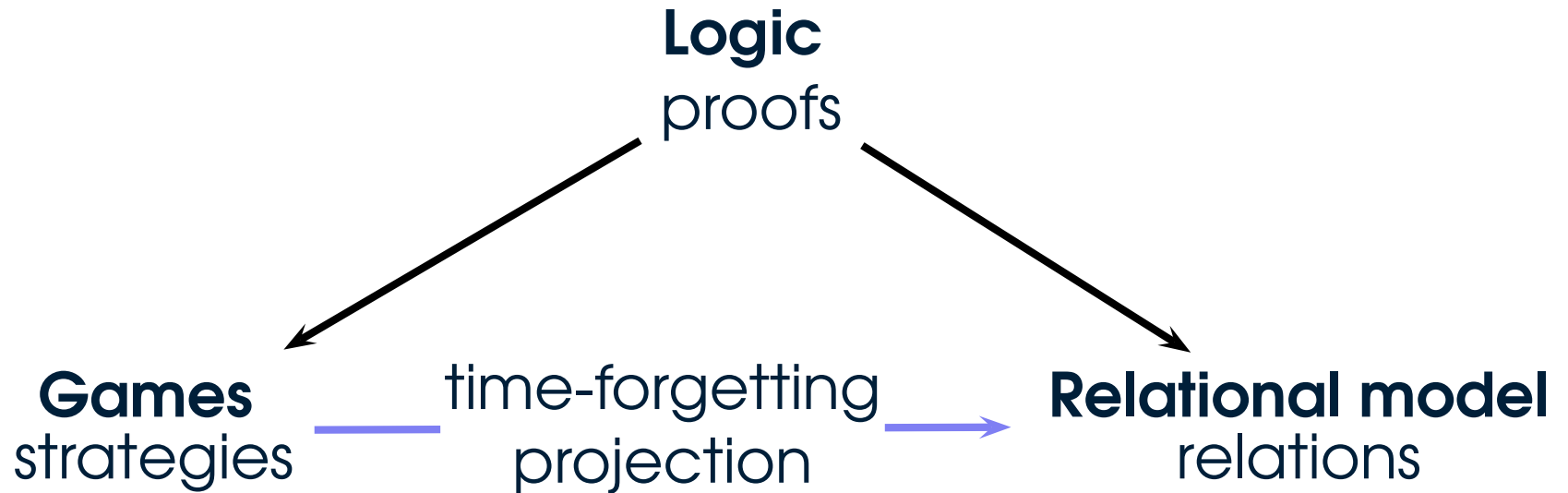
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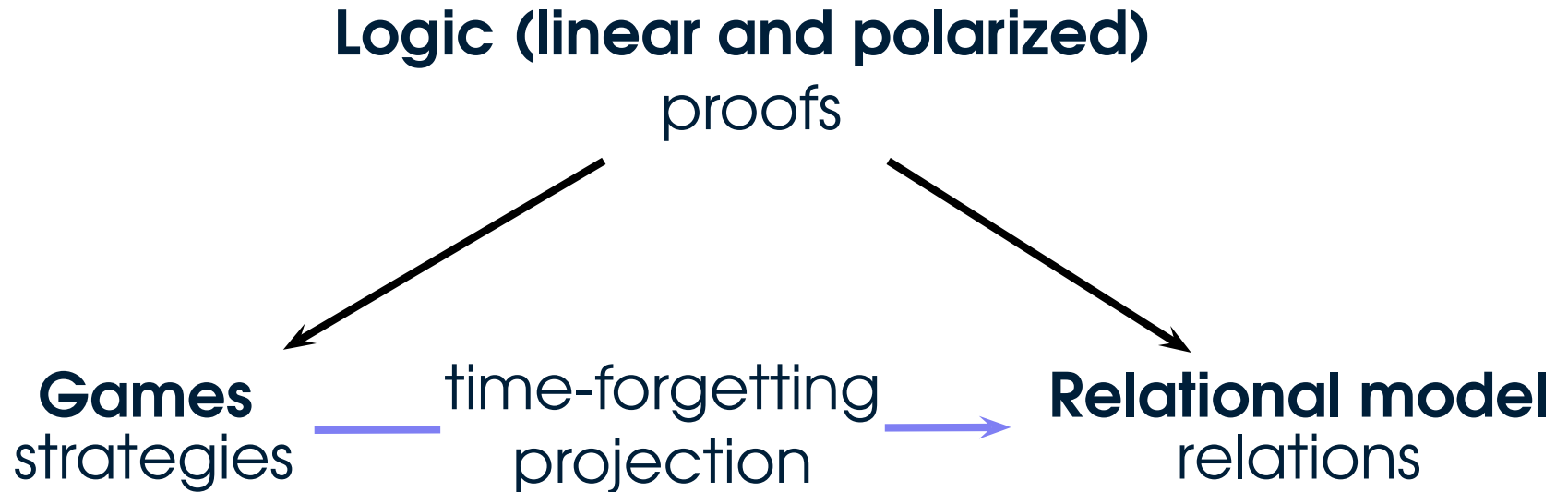
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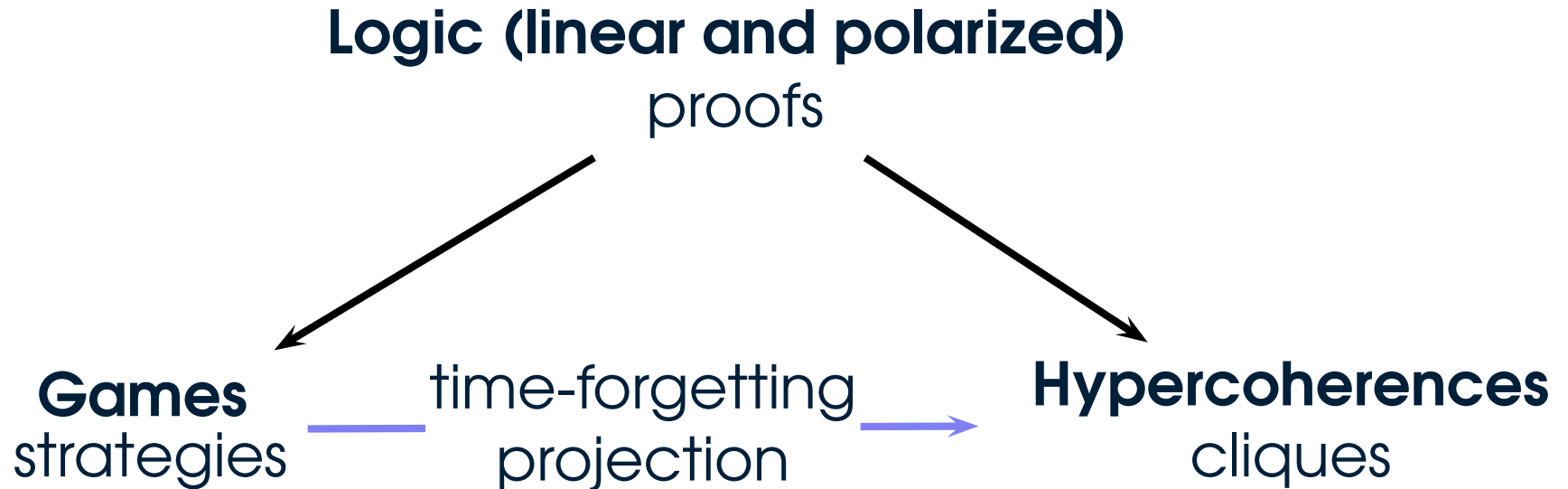
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- Here: commutation by use of polarized bordered games for the linear subsystem of linear logic with polarities (MALLPoI).

Linear logic with polarities (Laurent)

positives $P := 0 \mid 1 \mid \alpha^\perp \mid P \oplus P \mid P \otimes P \mid !N$

negatives $N := \top \mid \perp \mid \alpha \mid N \& N \mid N \wp N \mid ?P$

$$\frac{}{\vdash \alpha, \alpha^\perp} \text{ (ax.)} \quad \frac{}{\vdash 1} \text{ (one)} \quad \frac{\vdash \Gamma}{\vdash \perp, \Gamma} \text{ (bot)} \quad \frac{}{\vdash \top, \Gamma} \text{ (top)}$$

$$\frac{\vdash \Gamma, N \quad \vdash \Gamma, N'}{\vdash \Gamma, N \& N'} \text{ (with)} \quad \frac{\vdash \mathcal{N}, P_i \quad (i = 1, 2)}{\vdash \mathcal{N}, P_1 \oplus P_2} \text{ (plus)}$$

$$\frac{\vdash N, N', \Gamma}{\vdash N \wp N', \Gamma} \text{ (par)} \quad \frac{\vdash \mathcal{N}, P \quad \vdash \mathcal{N}', P'}{\vdash \mathcal{N}, \mathcal{N}', P \otimes P'} \text{ (tens.)} \quad \frac{\vdash \mathcal{N}, P}{\vdash \mathcal{N}, ?P} \text{ (der.)}$$

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$$\frac{\vdash \mathcal{N}, \mathcal{N}^\perp \quad \vdash \mathcal{N}, \Gamma}{\vdash \mathcal{N}, \Gamma} \text{ (cut)}$$

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Polarized bordered games

Games:

positive	$A = (+, A^o, A^p, S_A)$	$S_A \subseteq (A^p \cdot A^o)^*$
negative	$A = (-, A^o, A^p, S_A)$	$S_A \subseteq (A^o \cdot A^p)^*$

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Strategy $\sigma \subseteq S_A$ s.t.:

$$s, s' \in \sigma, s \neq s' \implies \text{length}(s \wedge s') \text{ is } \begin{cases} \text{even, if } A < 0 \\ \text{odd, if } A > 0 \end{cases}$$

Interleaving of words

➤ $A \cap B = \emptyset, u \in A^*, v \in B^*$

$$u \bullet_{A,B} v = \{w \mid w \upharpoonright A = u, w \upharpoonright B = v\}$$

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$$a \cdot u \cdot a' \otimes_{A,B} b \cdot v \cdot b' = \{(a, b) \cdot w \cdot (a', b') \mid w \in u \odot_{A,B} v\}$$

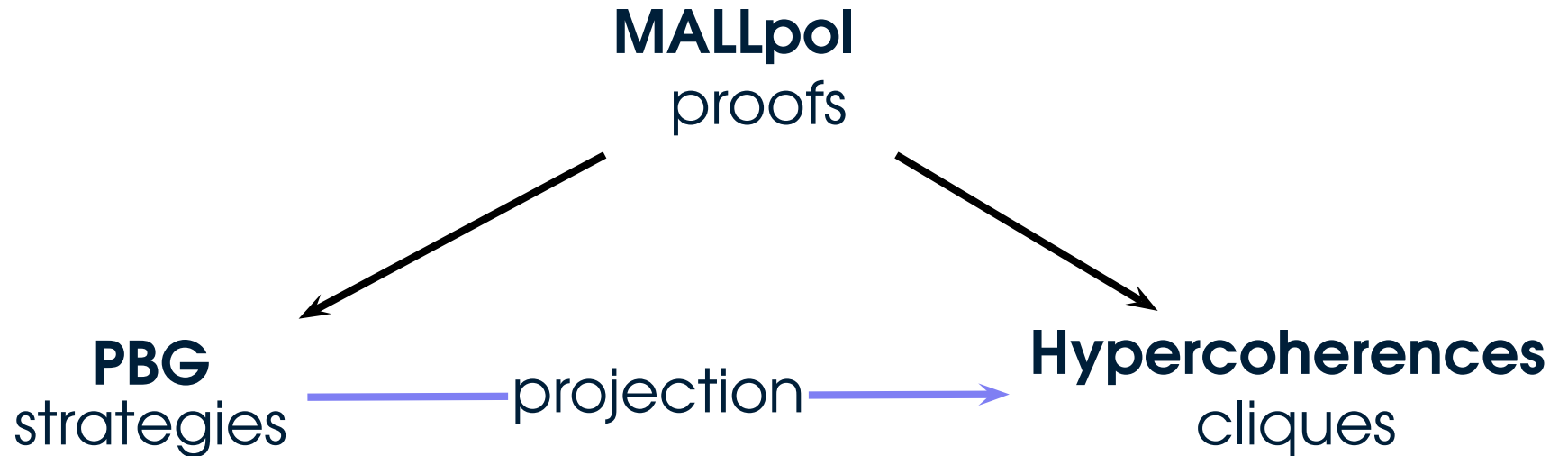
Interpretation of MALLpol/LLpol

- Follows the pattern of Laurent's polarized games 2002
 - Orthogonal: exchanges opponant and proponent
 - Additives: disjoint unions ($S_{A\oplus B} = S_A \uplus S_B$)
 - Multiplicatives: $S_{A\otimes B} = \{s \otimes_{A,B} s' \mid s \in S_A, s \in S_B\}$
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 - exponentials: $!N = \downarrow \#N$ (some choice for #)

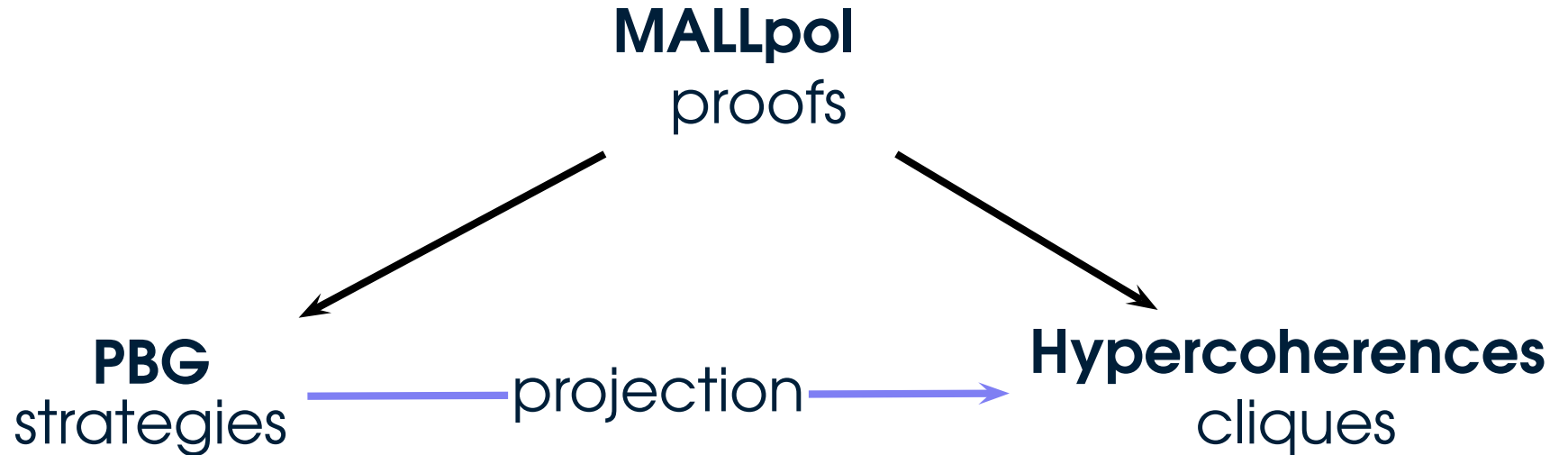
Projection



$$p([A]_{\text{PBG}}) = [A]_{\text{Rel}} = |[A]_{\text{Hc}}|$$

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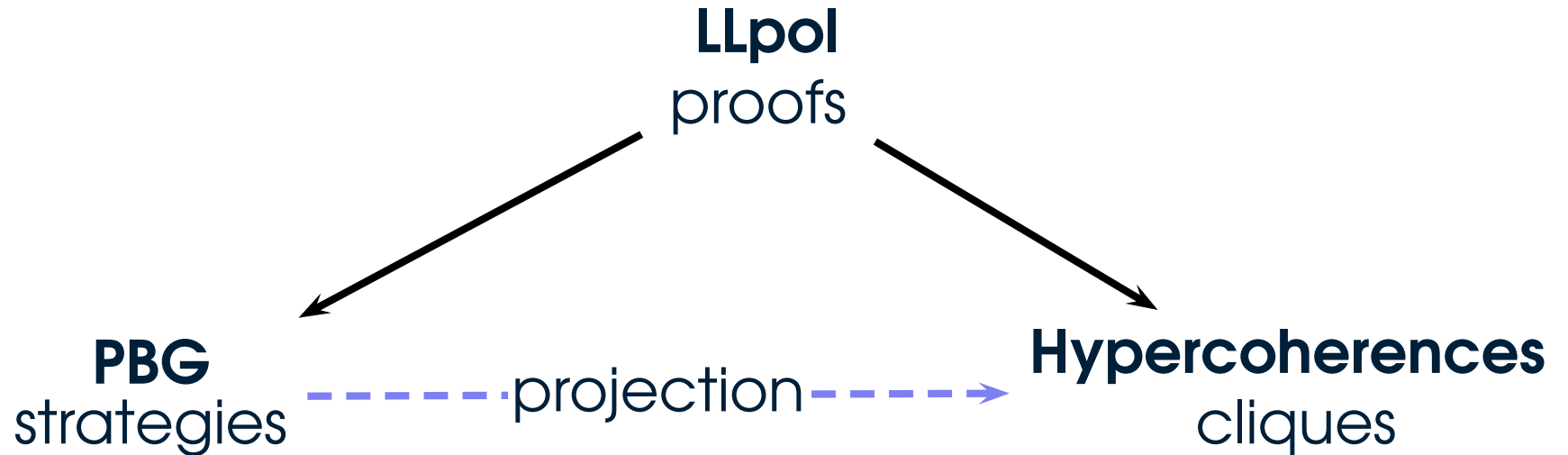
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Reversibility:

$$\text{reverse}(S_A) \cong S_A$$

$$\text{reverse}([\pi]_{\text{PBG}}) = [\pi]_{\text{PBG}}$$

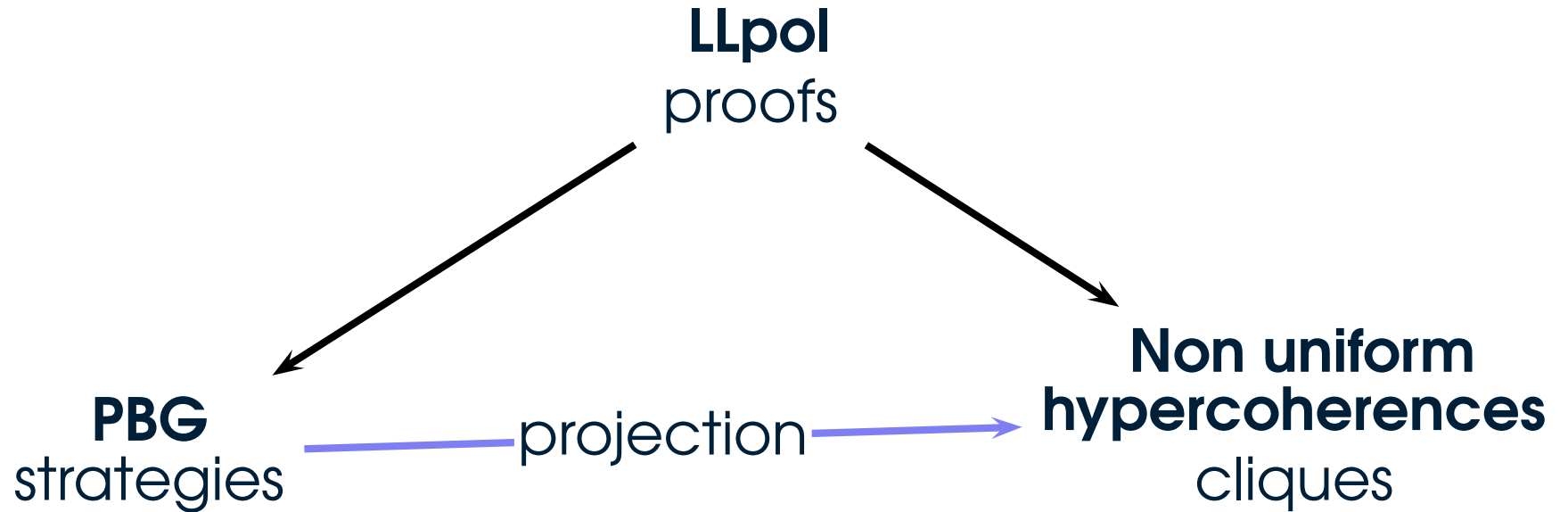
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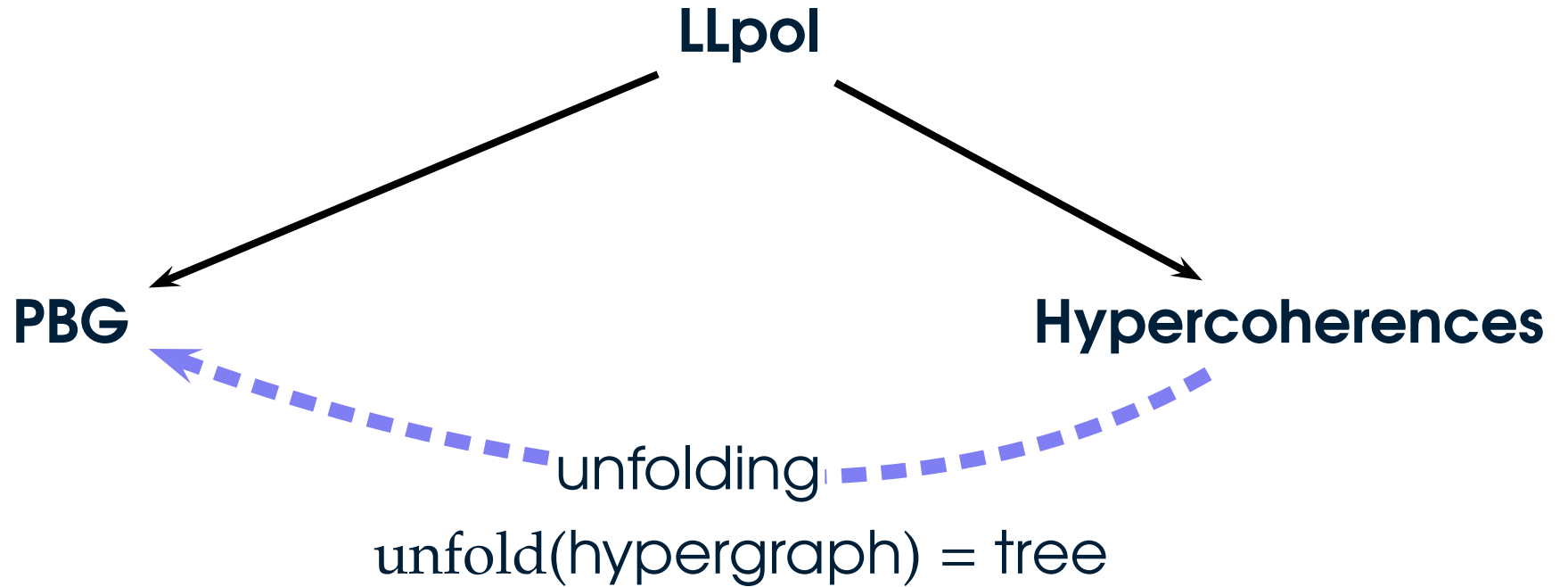
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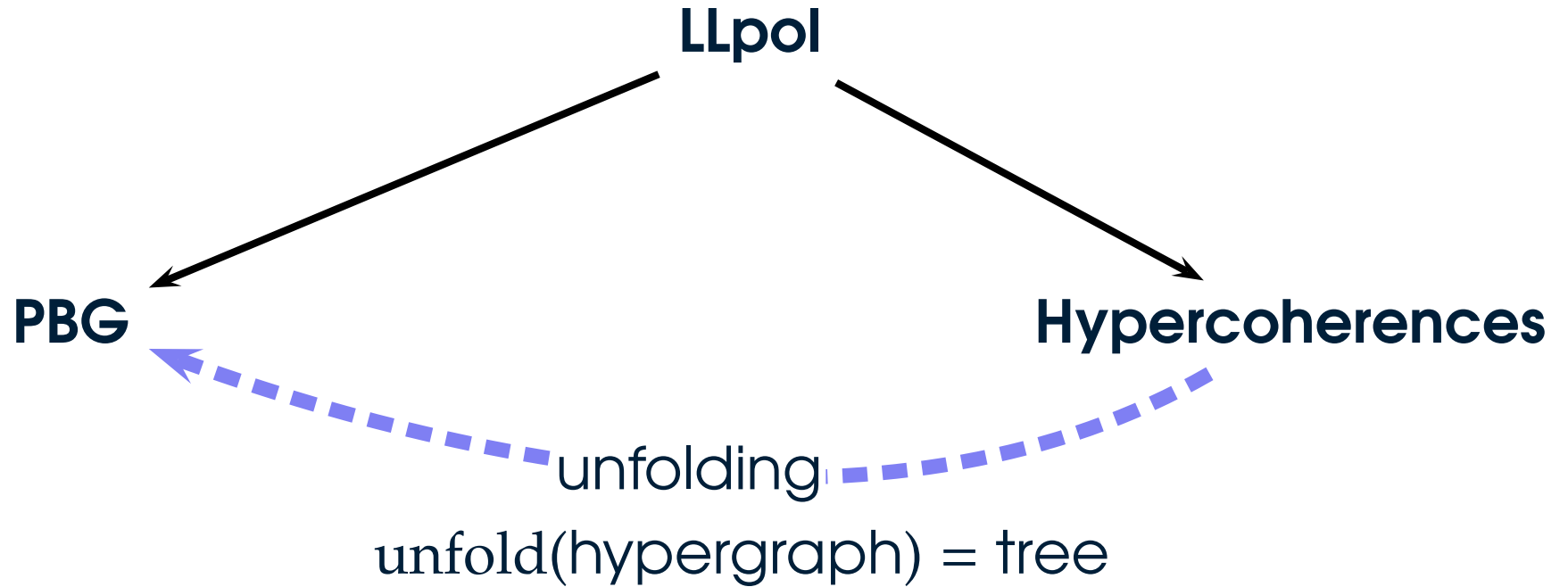
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Hypercoherences unfolding

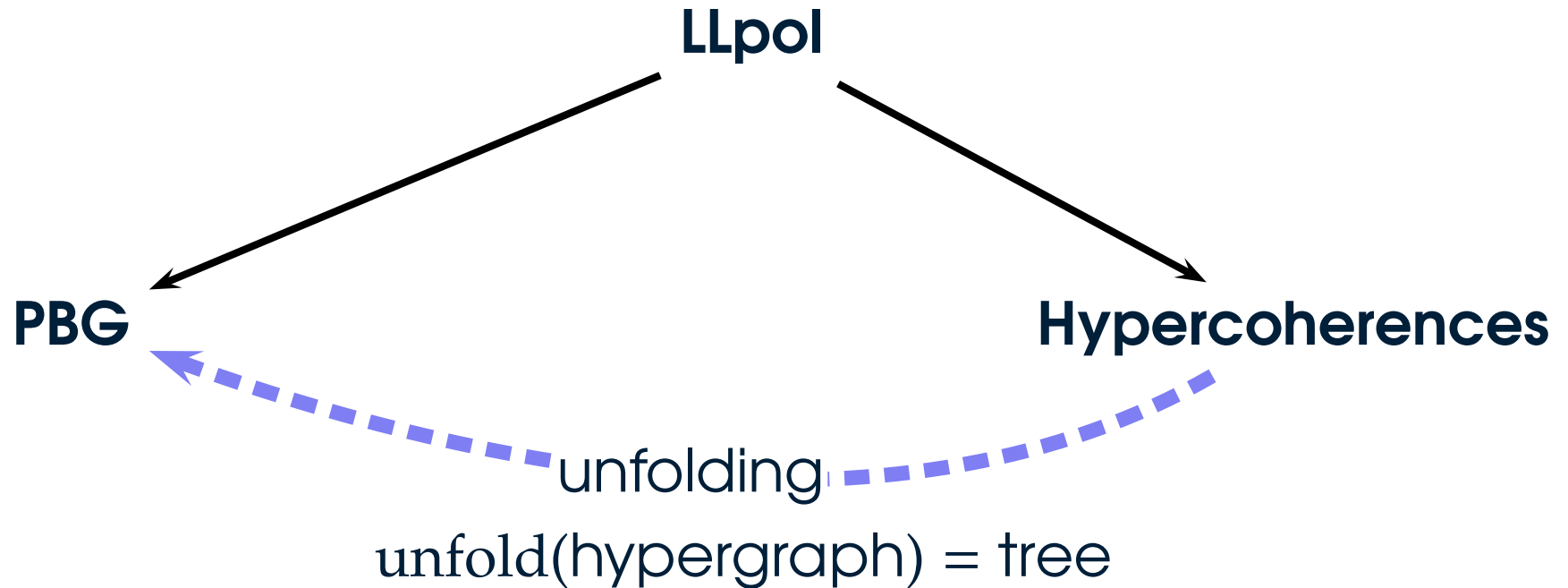


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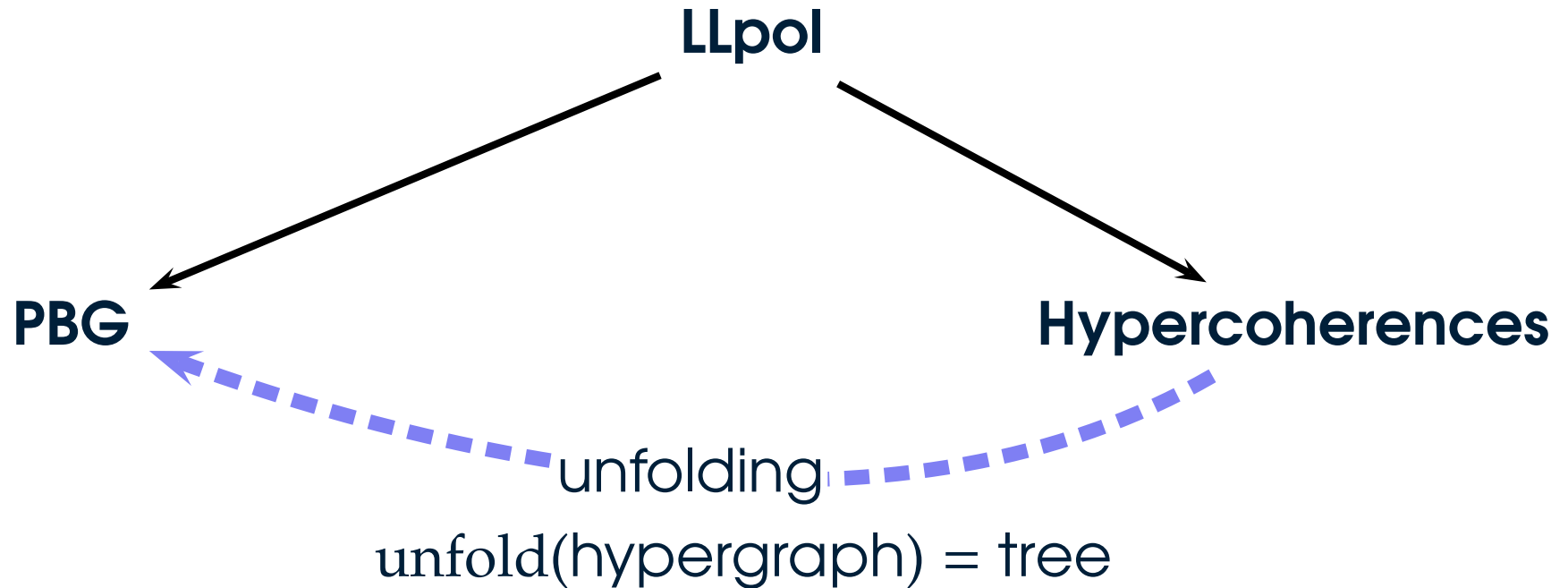
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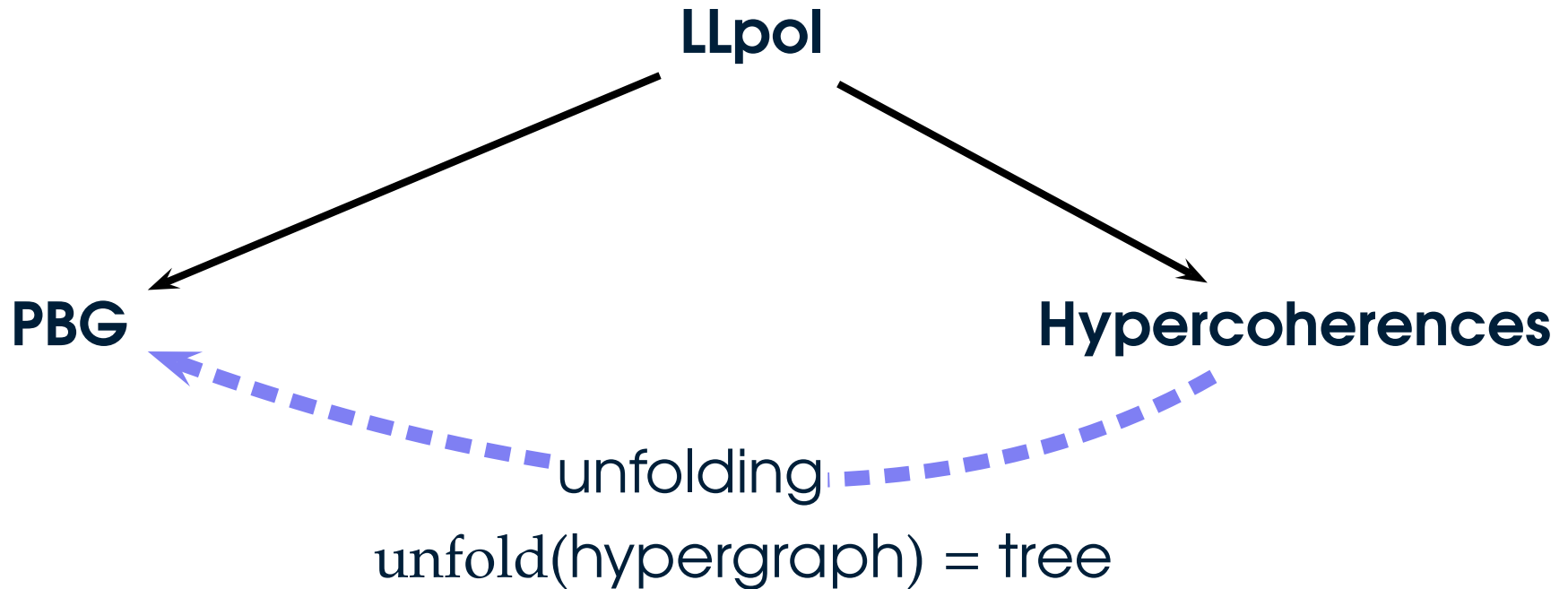
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 - $\text{unfold}(X \otimes Y) \cong \text{unfold}(X) \otimes \text{unfold}(Y)$
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 $\text{unfold}(!X) \cong \downarrow \#_s \text{unfold}(X)$

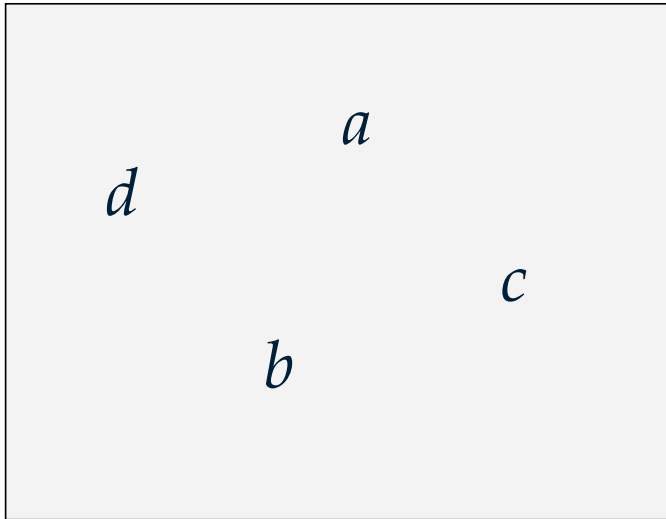
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- Fails in general.

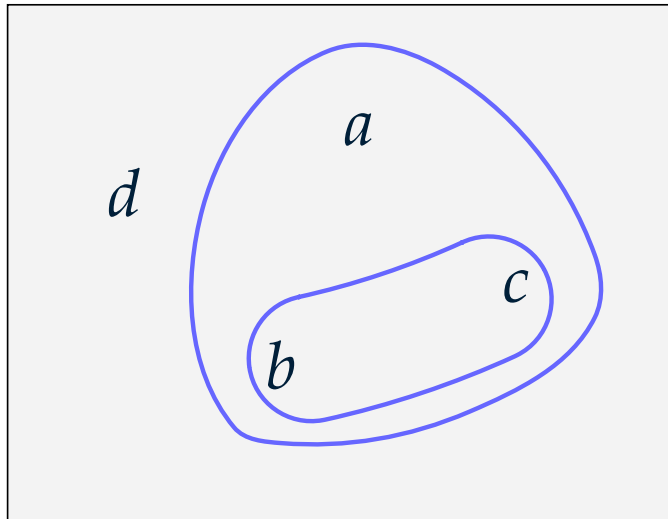
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hypercoherence



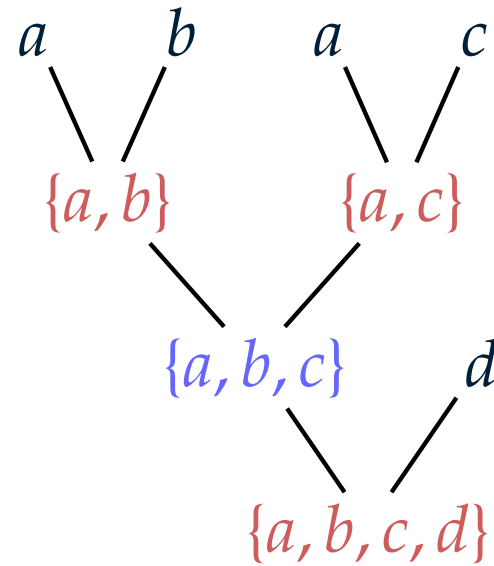
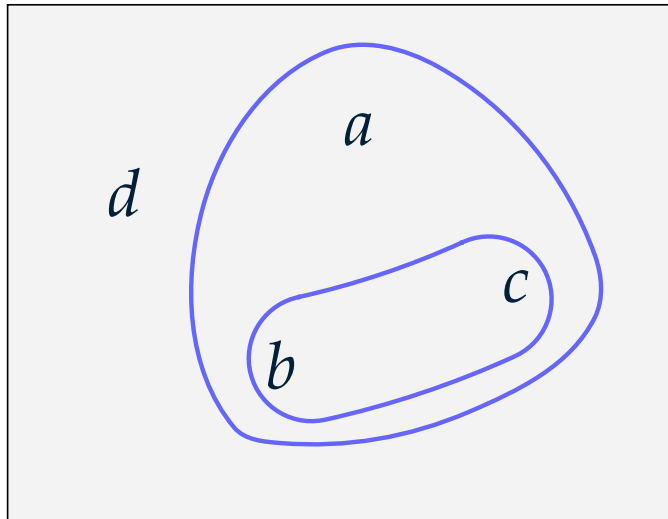
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